

Binary Search



Phone book

- Like linear search, binary search finds whether a certain item (the key) is in an array.
- Binary search only works on ***sorted*** arrays.
 - Binary search takes advantage of the array being sorted to make the search much faster.

key = 33

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
6	13	14	25	33	43	51	53	64	72	84	93	95	96	97

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6	13	14	25	33	43	51	53	64	72	84	93	95	96	97

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key = 33

Found! (return 4)

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
6	13	14	25	33	43	51	53	64	72	84	93	95	96	97

- We are given
 - an array (called **array**) of size n , indexed from 0 to $n-1$
 - an integer **key** to look for in the array
 - an integer **low** that is the lowest index in the array that could contain the key
 - an integer **high** that is the highest index in the array that could contain the key
- If **low** > **high**, then the item is not found (return -1)
- Compute the **middle** index in the array.
- If the item at the **middle** index is the key, return that index.
- If the item at the **middle** index is greater than the **key**, repeat from step 2, on the left sub-array.
- If the item at the **middle** index is less than the **key**, repeat from step 2 on the right sub-array.

- Three variables that do most of the work:
 - **low**: the smallest index that could possibly contain the key.
 - **high**: the largest index that could possibly contain the key.
 - **mid**: the midpoint of the two indices.

- If $low > high$, we know the item is not found (stop).
- If $array[mid] == key$, item is found (stop).
- If $array[mid] > key$, repeat algorithm with only the left half of the array.
- If $array[mid] < key$, repeat algorithm with only the right half of the array.

key = 33



0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
6	13	14	25	33	43	51	53	64	72	84	93	95	96	97

low = 0

high = 14

mid =

key = 33

low ↓								mid ↓							high ↓
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	
6	13	14	25	33	43	51	53	64	72	84	93	95	96	97	

low = 0
high = 14
mid = 7

key = 33

low			mid					high						
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
6	13	14	25	33	43	51	53	64	72	84	93	95	96	97

low = 0
high = 14
mid = 7

key = 33



0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
6	13	14	25	33	43	51	53	64	72	84	93	95	96	97

low = 0

high = 6

mid =

key = 33

low			mid		high									
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
6	13	14	25	33	43	51	53	64	72	84	93	95	96	97

low = 0

high = 6

mid = 3

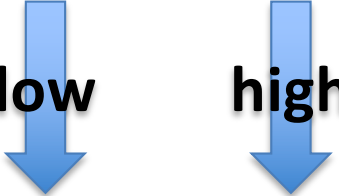
key = 33

low			mid		high									
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
6	13	14	25	33	43	51	53	64	72	84	93	95	96	97

low = 0
high = 6
mid = 3

key = 33

low **high**



0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
6	13	14	25	33	43	51	53	64	72	84	93	95	96	97

low = 4

high = 6

mid =

key = 33

low mid high



0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
6	13	14	25	33	43	51	53	64	72	84	93	95	96	97



low = 4

high = 6

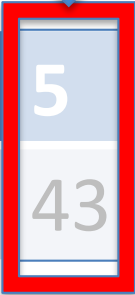
mid = 5

key = 33

low mid high



0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
6	13	14	25	33	43	51	53	64	72	84	93	95	96	97



low = 4

high = 6

mid = 5

key = 33



0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
6	13	14	25	33	43	51	53	64	72	84	93	95	96	97

low = 4
high = 4
mid =

key = 33
Found!

high
mid
low



0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
6	13	14	25	33	43	51	53	64	72	84	93	95	96	97

low = 4
high = 4
mid = 4

- If $\text{low} > \text{high}$, we know the item is not found (stop).
- If $\text{array}[\text{mid}] == \text{key}$, item is found (stop).
- If $\text{array}[\text{mid}] > \text{key}$, repeat algorithm ***with only the left half of the array.***
 - How do we change low & high?
- If $\text{array}[\text{mid}] < \text{key}$, repeat algorithm ***with only the right half of the array.***
 - How do we change low & high?

- If $low > high$, we know the item is not found (stop).
- If $array[mid] == key$, item is found (stop).
- If $array[mid] > key$, repeat algorithm ***with only the left half of the array.***
 - How do we change low & high?
 - **$high = mid - 1$**
- If $array[mid] < key$, repeat algorithm ***with only the right half of the array.***
 - How do we change low & high?
 - **$low = mid + 1$**

Recursive formulation

- Function: `binarySearch(array, key, low, high)`
- Base cases:
 - If key is found: Return position found.
 - If $low > high$: Return -1 (indicating not found).
- Recursive cases:
 - If $array[mid] > key$: `binarySearch(array, key, low, mid - 1)`
 - If $array[mid] < key$: `binarySearch(array, key, mid + 1, high)`