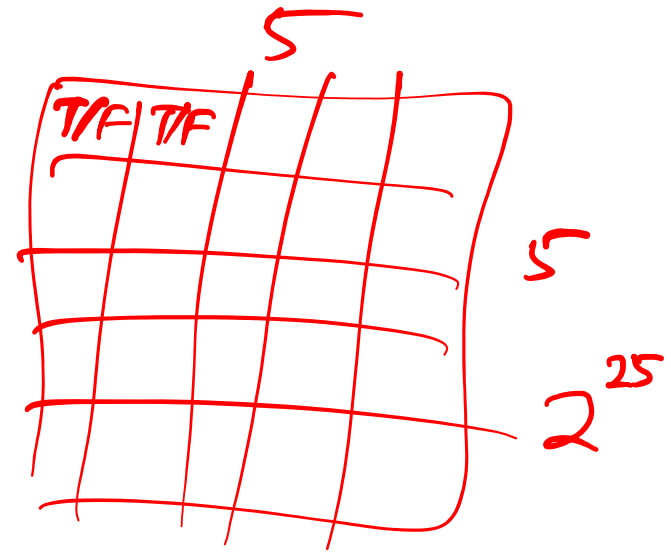
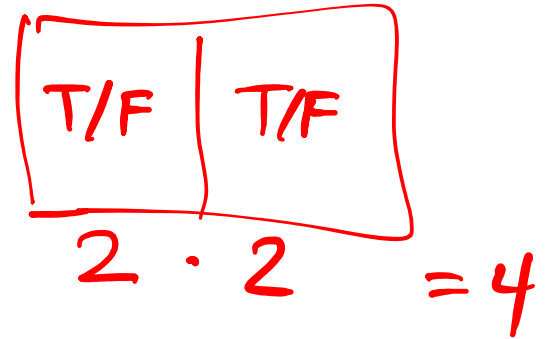


Formulate Roomba problem

location of Roomba = 2^x

clean/dirty for
each room = 4

8



Recap

- What things do we need to define in order to formulate a problem as a search problem?

Env

State

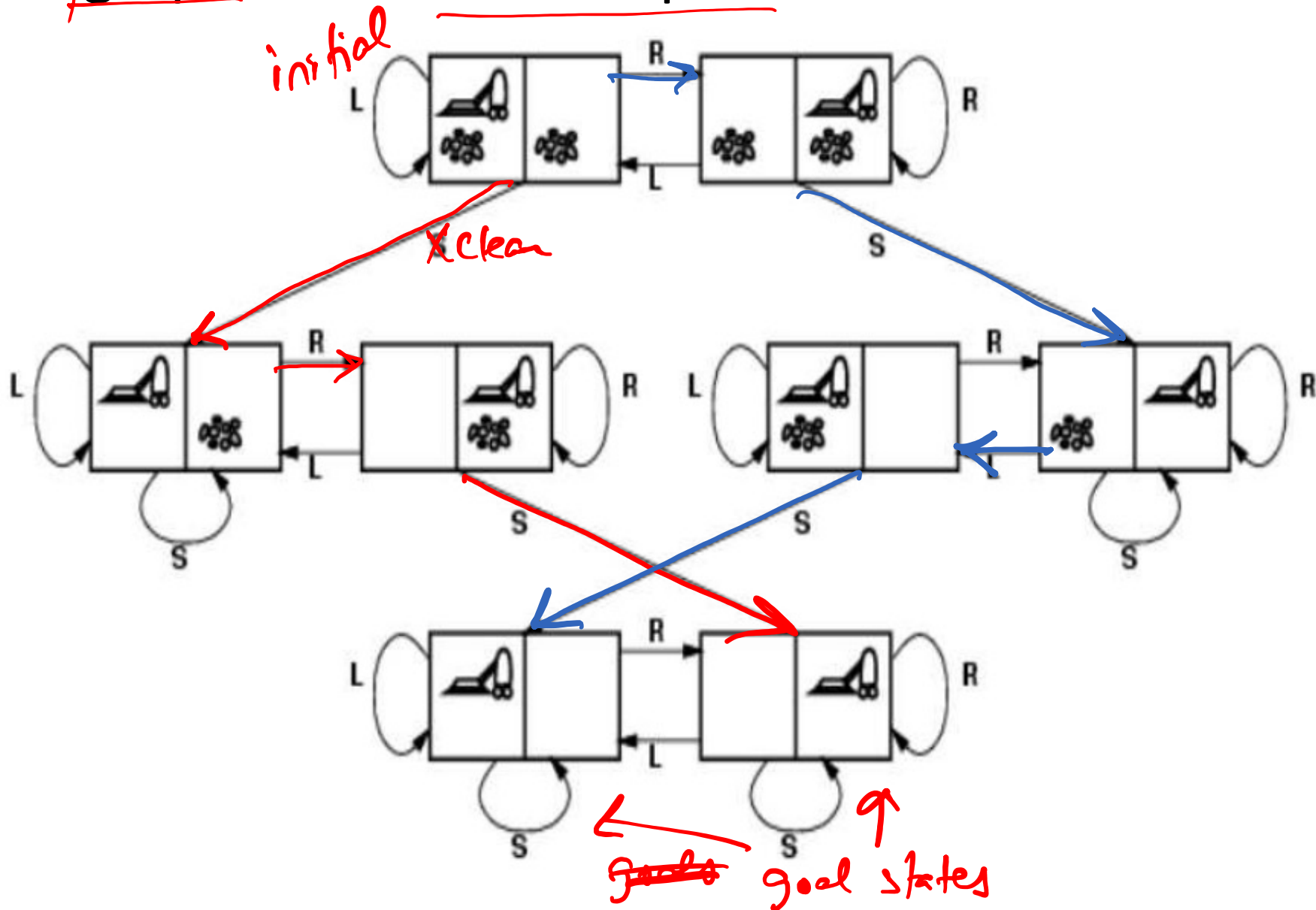
Initial state

Actions

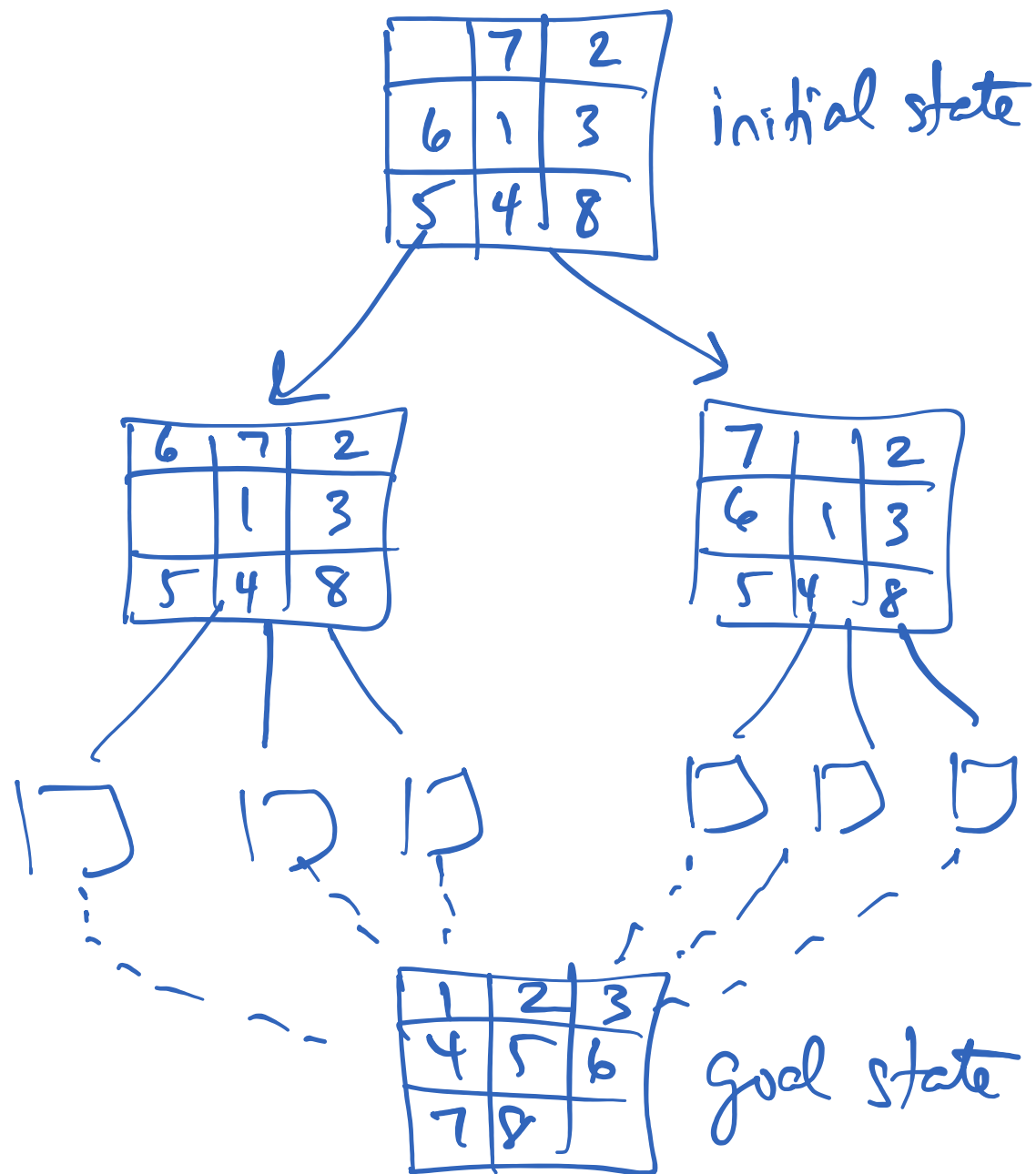
Goal state

Cost

- Always a good idea to try to visualize the graph of the search space.



8-puzzle/sliding block puzzle

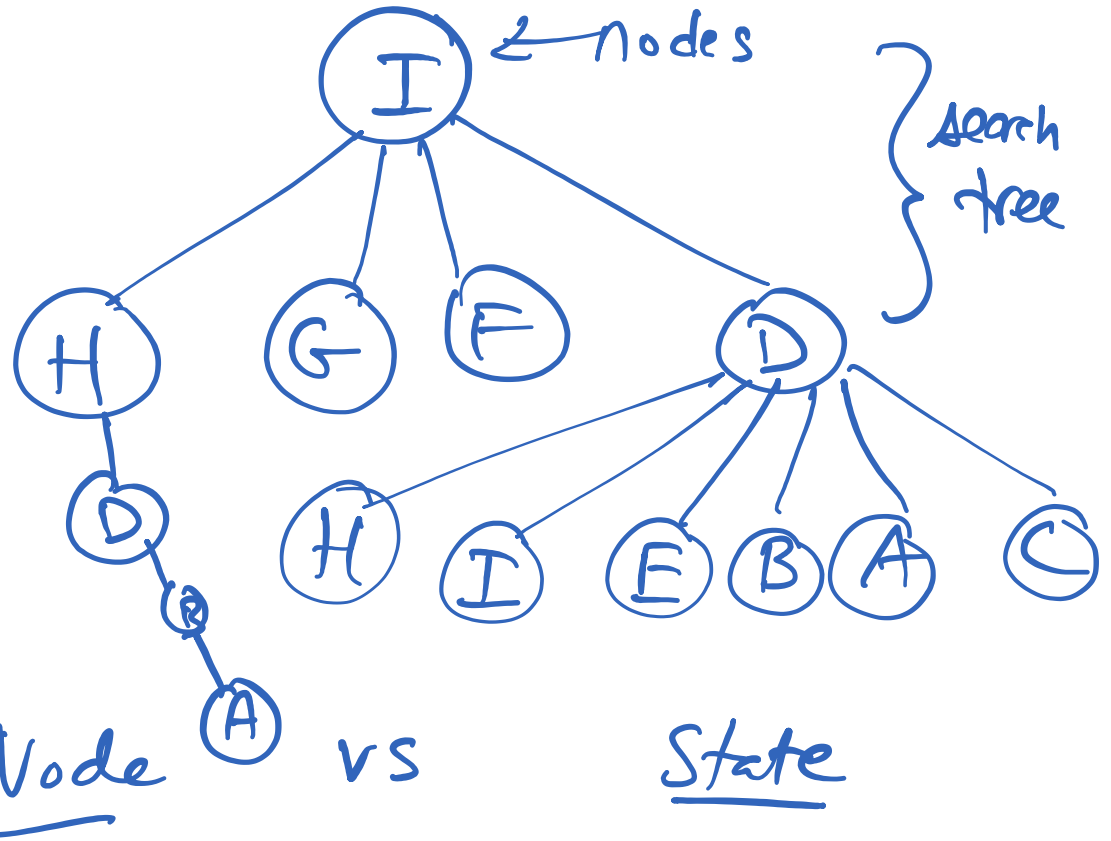
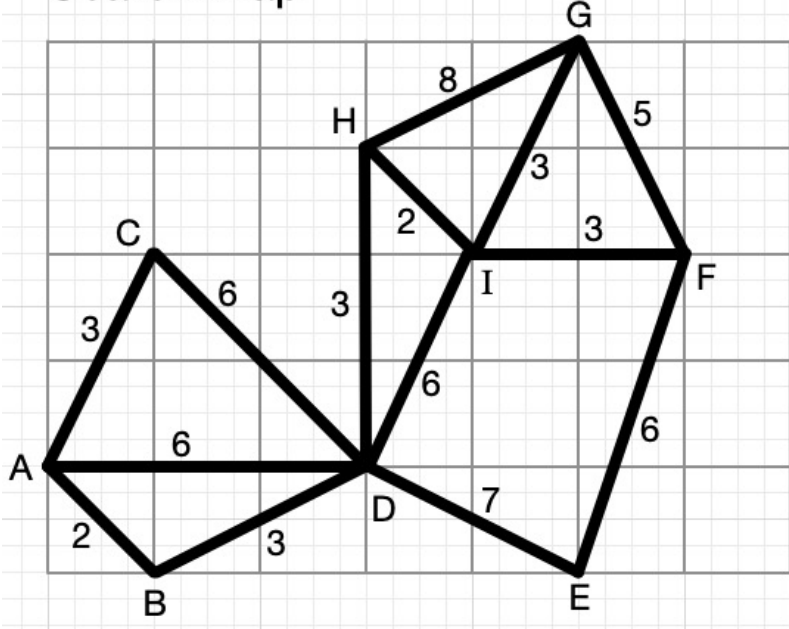


Generic search algorithms (3.3)

- All search algorithms work in essentially the same manner:
- Start with initial state.
- Generate all possible successor states (a.k.a. "expanding a node.") *"next states"*
- Pick a new node to expand. *← Differentiates each algorithm.*
- Continue until we find a goal state.

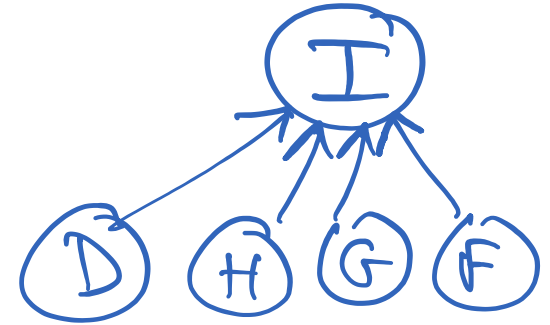
Search Map

$I \rightarrow A$



- There are two simultaneous graph-like structures used in search algorithms:
 - (1) Tree or graph of the underlying state space.
almost always
 - (2) Tree maintaining the record of the current search in progress (the ***search tree***).
- (1) does not depend on the current search being run.
- (1) is sometimes not even stored in memory (too big!)
- (2) always depends on the current search, and is always stored in memory. It is created on the fly during the running of the search algorithm.

Search tree



- A node n of the search tree stores:
 - a state (of the state space)
 - a pointer ^{reference} to the state's parent node (usually)
 - the action that got you from the parent to n (sometimes)
 - the path cost $g(n)$: cost of the path so far from the initial state to n .

Generic search algorithms

(all based off of "best-first search")

- **Frontier**: a data structure storing the collection of nodes that are available to be examined next in the algorithm.
 - Often represented as a stack, queue, or priority queue.
- **Reached**: a map from nodes to states. Keeps track of which states have been examined already.
 - Often stored using a data structure that enables quick look-up for membership tests.

key value

hash table
~~BST~~




How do you evaluate a search algorithm?

- **Completeness** — Does the algorithm always find a solution if one exists? ↑ lowest cost
- **Optimality** — Does the algorithm find the best solution?
- **Time complexity** — ~~⊗~~ $O(n)$
- **Space complexity** — measure memory usage — $O(n)$

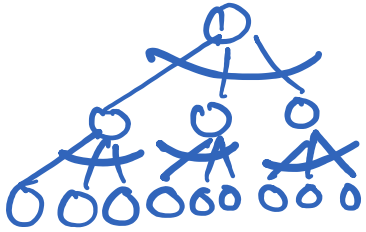
Uninformed search methods

- These methods have no information about which nodes are on promising paths to a solution.
- Also called: *blind search*

Uninformed Search algorithms

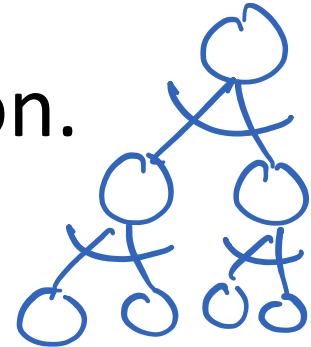
-  Breadth-first search
-  Uniform-cost search
-  Depth-first search

```
function BREADTH-FIRST-SEARCH(problem) returns a solution node or failure  
  node  $\leftarrow$  NODE(problem.INITIAL)  
  if problem.IS-GOAL(node.STATE) then return node  
  frontier  $\leftarrow$  a FIFO queue, with node as an element  
  reached  $\leftarrow$  {problem.INITIAL}  
  while not IS-EMPTY(frontier) do  
    node  $\leftarrow$  POP(frontier)  
    for each child in EXPAND(problem, node) do  
      s  $\leftarrow$  child.STATE  
      if problem.IS-GOAL(s) then return child  
      if s is not in reached then  
        add s to reached  
        add child to frontier  
  return failure
```



Breadth-first search

- Choose shallowest node for expansion.
- Data structure for frontier?
 - Queue (regular)



- Complete? Optimal? Time? Space? — $O(b^d)$

Yes

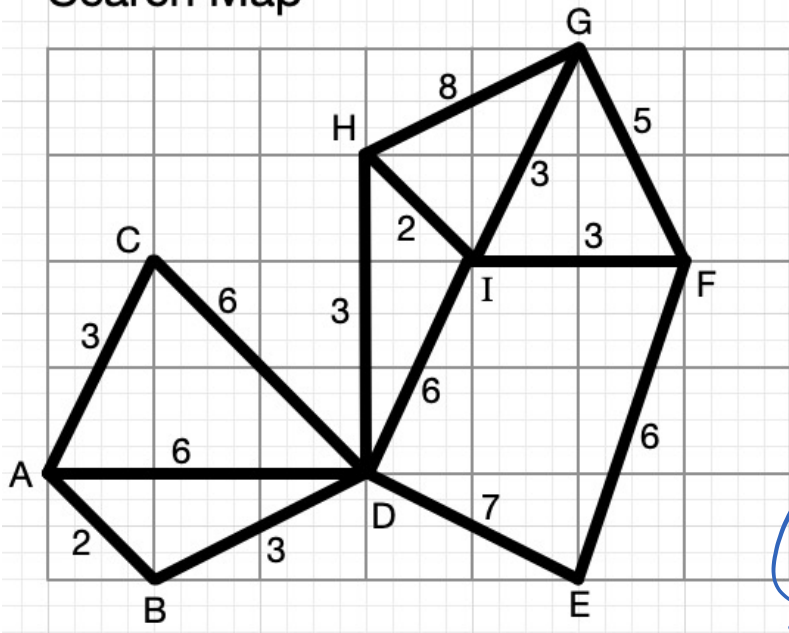
No, unless...
all costs
are the
same

(b)
— branching factor: max #
of actions from a state
— depth of the solution in
the search tree (d)

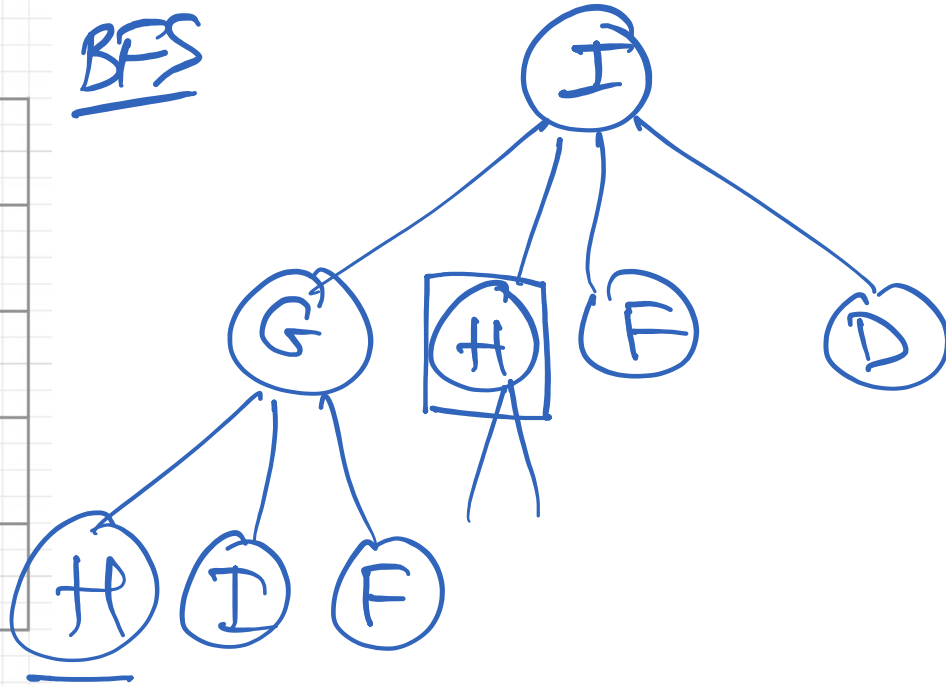
$$O(b^d)$$

"n"
"Size of the input"

Search Map



BFS



Depth-first search (DFS)

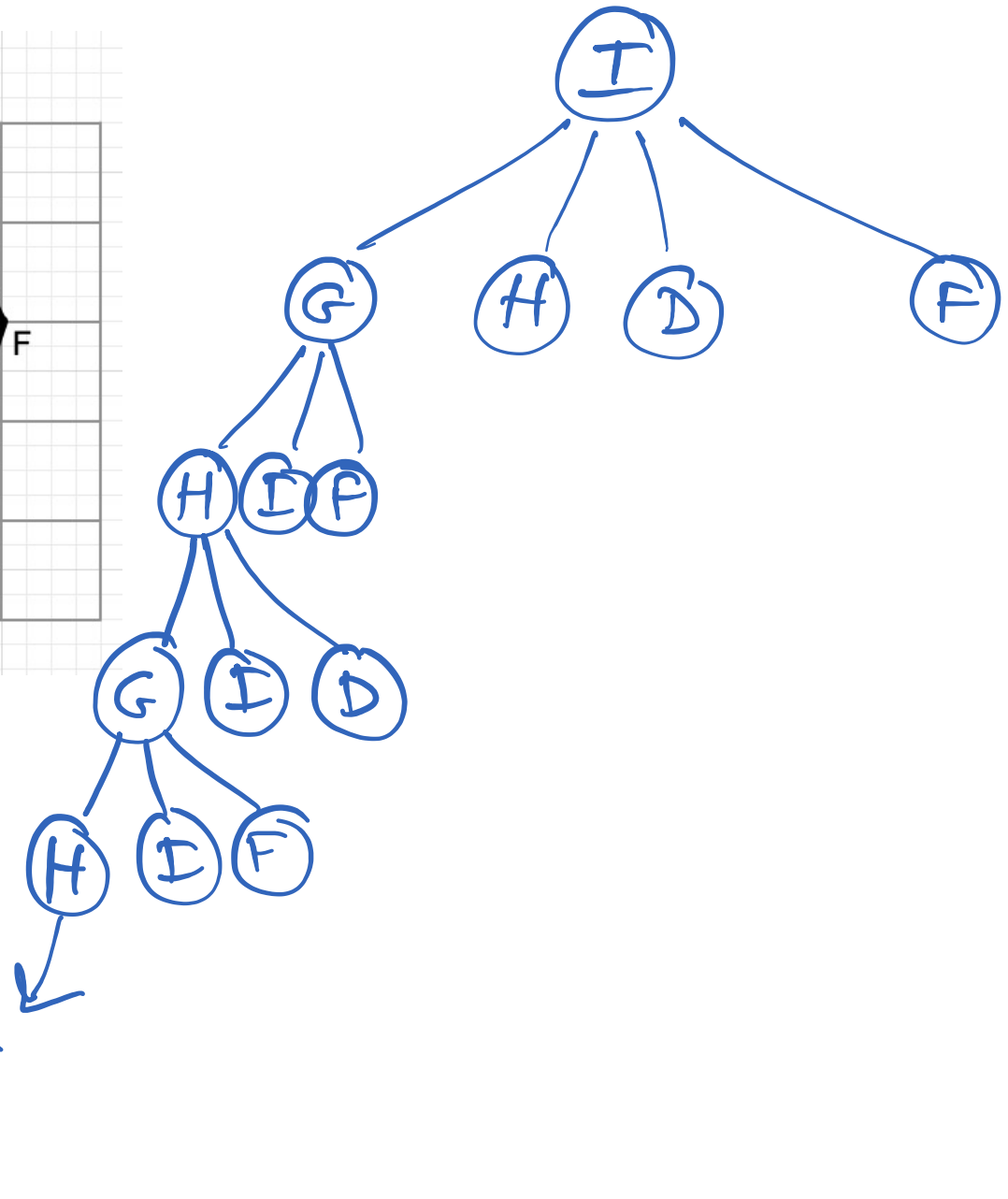
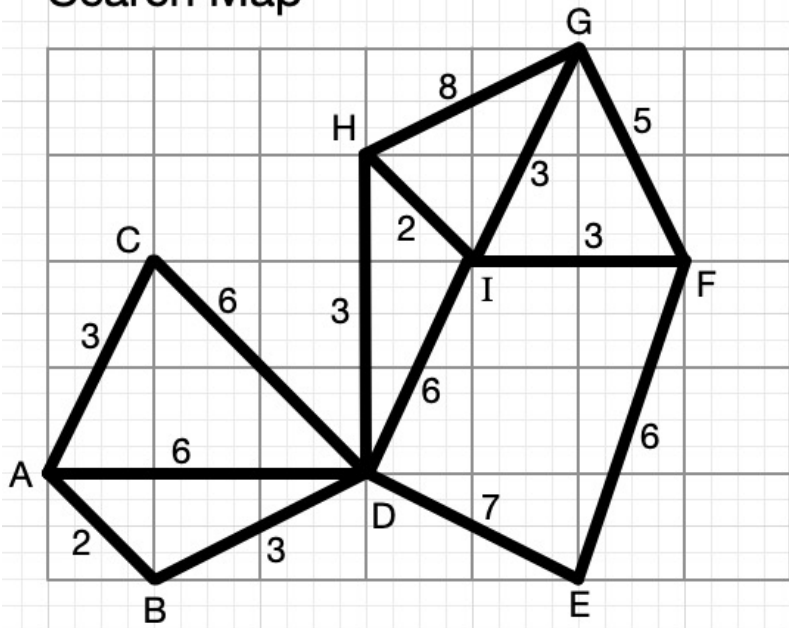
- Choose deepest node to expand.
- Data structure for frontier?
 - Stack (or just use recursion)
- Complete? Optimal? Time? Space?

↳ No, unless
you prevent
loops.

(Yes, if you
do)

No

Search Map



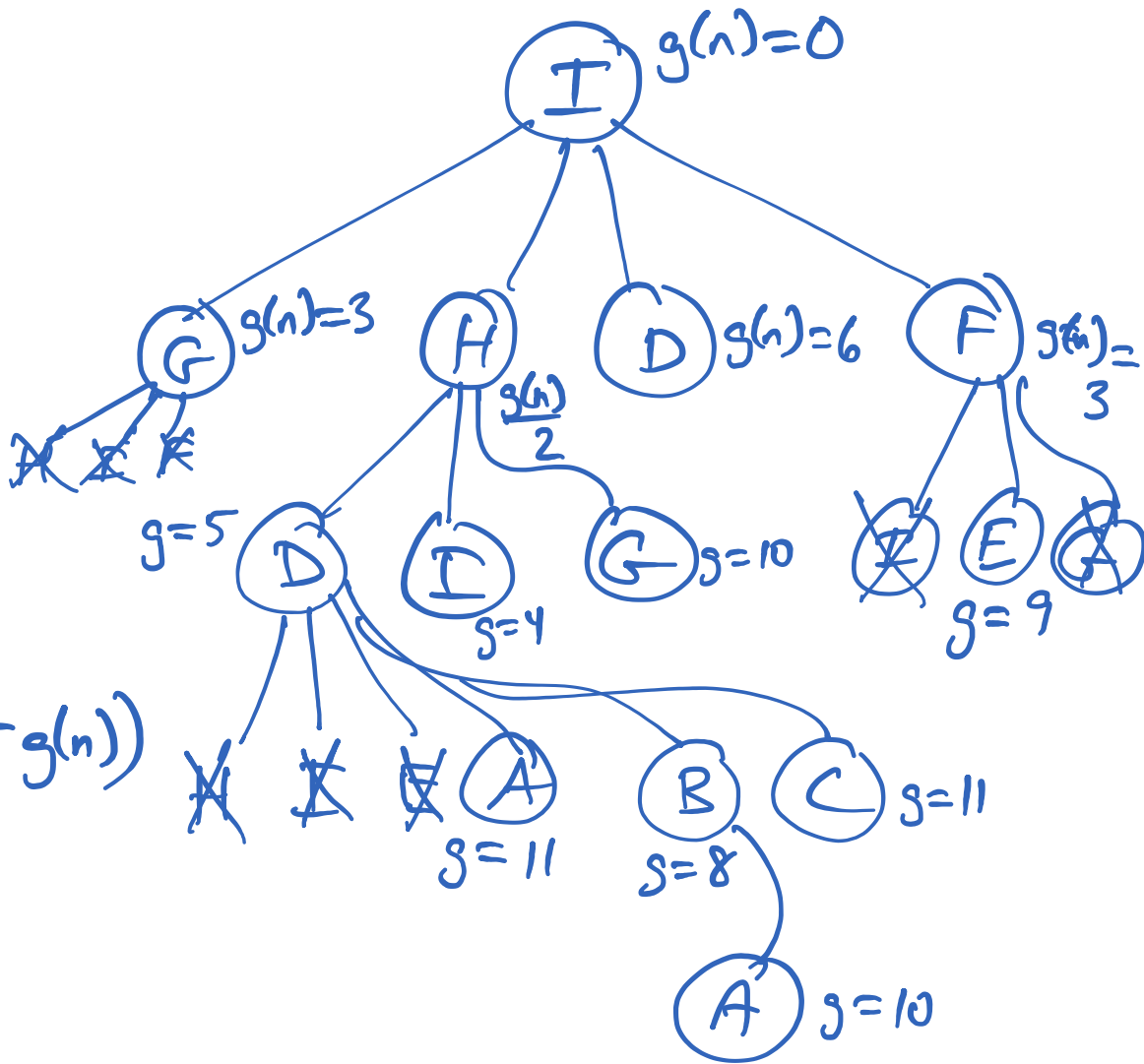
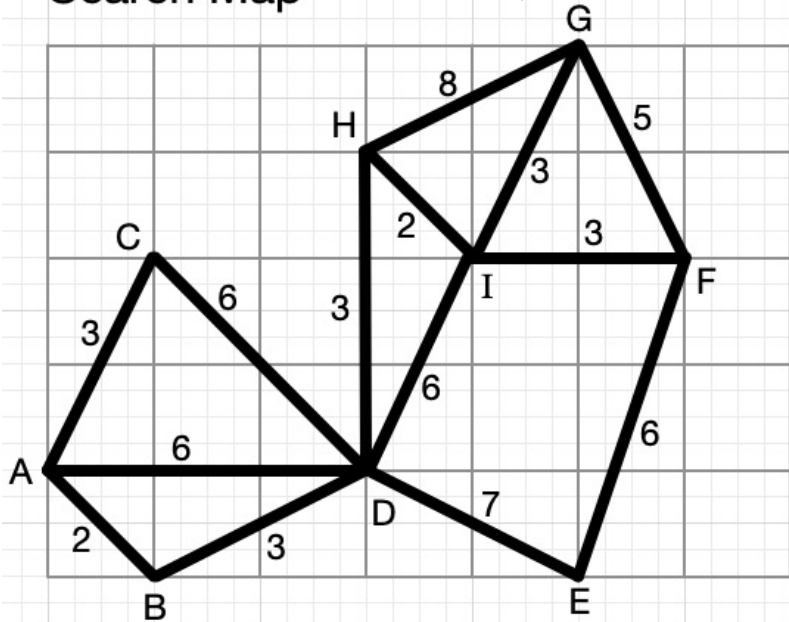
Dijkstra = Uniform-cost search

total cost so far from initial state to n

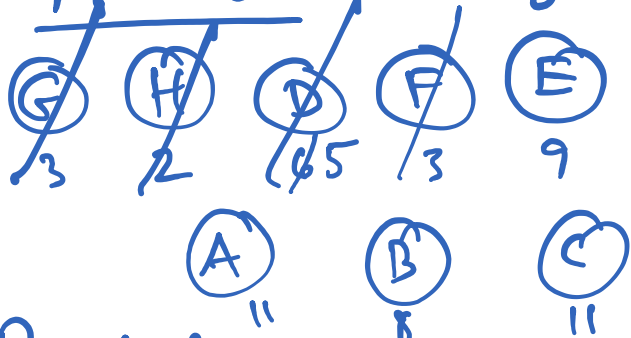
- Choose node with lowest path cost $g(n)$ for expansion.
- Data structure for frontier?
 - Priority queue
- Suppose we come upon the same state twice. Do we re-add to the frontier?
 - Yes, if lower path cost.

Search Map

I → A



Frontier (sorted by lowest $g(n)$)



Reached

I G H ~~D~~ F E

I H D B A = 10